Ranking comment sorting policies in online debates

Anthony P. Young a,*, Sagar Joglekar b, Gioia Boschi b and Nishanth Sastry a, c

a Department of Informatics, King’s College London, Bush House, Strand Campus, 30 Aldwych, London, WC2B 4BG, United Kingdom
E-mails: peter.young@kcl.ac.uk, sagar.joglekar@kcl.ac.uk, nishanth.sastry@kcl.ac.uk

b Department of Mathematics, King’s College London, Strand, London, WC2R 2LS, United Kingdom
E-mail: gioia.boschi@kcl.ac.uk

c Department of Computer Science, University of Surrey, Guildford, Surrey, GU2 7XH, United Kingdom
E-mail: n.sastry@surrey.ac.uk

Abstract. Online debates typically possess a large number of argumentative comments. Most readers who would like to see which comments are winning arguments often only read a part of the debate. Many platforms that host such debates allow for the comments to be sorted, say from the earliest to latest. How can argumentation theory be used to evaluate the effectiveness of such policies of sorting comments, in terms of the actually winning arguments displayed to a reader who may not have read the whole debate? We devise a pipeline that captures an online debate tree as a bipolar argumentation framework (BAF), which is sorted depending on the policy, giving a sequence of induced sub-BAFs representing how and how much of the debate has been read. Each sub-BAF has its own set of winning arguments, which can be quantitatively compared to the set of winning arguments of the whole BAF. We apply this pipeline to evaluate policies on Kialo debates, where it is shown that reading comments from most to least liked, on average, displays more winners than reading comments earliest first. Therefore, in Kialo, reading comments from most to least liked is on average more effective than reading from the earliest to the most recent.

Keywords: Online debates, comment sorting policies, bipolar argumentation frameworks, Kialo

1. Introduction

1.1. Background and motivations

The Internet has enabled people to share their views about many topics online, ranging from commenting on important news events1 to debating philosophical stances.2 Given that the number of Internet users grew eight-fold to over three billion worldwide between 2000 to 2016 [27], it is not surprising that such discussions have also grown both in the number of comments posted and the range of topics discussed [19]. Many online discussions now contain so many comments that it is unrealistic for a normal Internet user to read every single point being made; it is expected that a user interested in a given topic of debate would read only a fragment of the entire discussion. To help users, many online platforms that host

---

such discussions allow for the comments to be sorted according to some policy, such as displaying the comments from the most liked to the least liked.

Further, comments exchanged over such discussions are often argumentative in nature, and can be modelled with argumentation theory (e.g. [7,31]); this can resolve disagreements arising from such online debates by finding the winning arguments based on normative yet intuitive criteria. Suppose that a user would like to know which arguments have prevailed at a given moment, but does not have time to read every comment. Such a reader may think that a given argument has won when it has been rebutted by further comments which have not yet been read. The goal of this paper is to compare how various comment sorting policies can affect the number of actually winning arguments exposed to a reader, depending on how much of the debate has been read. A precise way to compare such policies will bring visibility to which arguments should win, thus helping readers to better navigate and understand large and argumentative online debates (e.g. [1,8,14,28,32,36]).

1.2. Overview of contributions

We apply argumentation theory, data mining and statistics to build a pipeline that compares comment sorting policies by measuring the number of actually winning arguments each policy displays to a reader who has only read a part of the debate. Intuitively, we mine online debates (e.g. such as [13,17,35]), and represent them as directed graphs (digraphs), whose nodes are the comments made during the debates and the edges denote which comment is replying to which other comment. As such debates begin with an initial comment, and each subsequent comment replies to exactly one other comment, these debates are trees, which is what we will restrict our analysis to in this paper, leaving other debate network topologies for future work (see Section 6).

Argumentation theory (e.g. [31]) is the branch of artificial intelligence concerned with the rational and transparent resolution of disagreements. Abstractly, arguments and their disagreements (called attacks) are represented as digraphs, called argumentation frameworks (AFs) [15]. However, in real debates, arguments can agree with each other as well as disagree. Subsequent theoretical developments have incorporated support between arguments, which can be interpreted as “agreement”. An AF where each directed edge is either an attack or a support is called a bipolar argumentation framework (BAF) [9,10]. There are principled ways of determining from various combinations of attacks and supports which arguments are winning. Therefore, subject to an appropriate representation of the debate one is interested in, resolving the disagreements is modelled by finding a subset of winning arguments.

To enable the application of argumentation theory, the pipeline makes two idealised assumptions: (1) each comment qualifies as a self-contained argument, and (2) each reply is either supporting or attacking. Of course, the vast majority of online debates are not so “clean”, but the pipeline does not preclude methods that allow for the identification of whether a comment qualifies as an argument (as opposed to, e.g. an insult), or whether a reply between comments is an attack or a support (e.g. [3,11]). Assuming that there is a relatively “clean” online debate, or that we can incorporate sophisticated data cleaning techniques to the pipeline, then we can treat the online debate formally as a BAF. Furthermore, if this debate is a tree, then the set of winning arguments is unique [15, Theorem 30].

To model the idea of a comment sorting policy, the pipeline linearly sorts all comments based on the direction of replies, such that if comment $b$ replies to comment $a$, then $a$ is read prior to $b$, while preserving the structure of each thread (see Section 3.1). However, when multiple comments $b_1, b_2, \ldots, b_n$ are replying to a given comment $a$, we will sort the $b_i$’s based on a given attribute. For this paper, we
will only consider the attributes \textit{likes} (represented by some number, see Section 5.2) and \textit{time} (represented e.g. by Unix time). Therefore, a comment sorting policy determines the order by which a reader is exposed to the arguments made in a debate with respect to an attribute.

Suppose the debate has $N \in \mathbb{N}^+$ comments ($N$ is a positive integer). The reader will read the first $n$ ($0 \leq n \leq N$) comments according to the policy, and for each such $n$, we will obtain an induced BAF of those comment-nodes ("induced", i.e. keeping all edges between those $n$ nodes). Each such induced BAF will have its own unique winning set of arguments $G_n$, which may or may not contain the set of \textit{actually winning} arguments $G_N$, i.e. the winning arguments after reading the whole debate. To quantitatively compare policies, the pipeline compares $G_n$ with $G_N$ for each $n$, which gives a sequence of numbers $J_0, J_1, J_2, \ldots, J_N$, where $J_n \in [0, 1]$ is the \textit{Jaccard coefficient} between the sets $G_n$ and $G_N$ (see Section 3). Intuitively, the larger $J_n$ is, the higher the proportion of actually winning arguments the policy exposes to a reader who has read $n$ comments. The pipeline calculates, for each policy, an $(N+1)$-tuple of $J_n$'s.

The pipeline then aggregates the $J_n$'s of each policy into a single number, such as taking the average (see Section 3.3), in order to compare which policy is more effective at displaying the actually winning arguments. Alternatively, one can fix $n$ and just compare $J_n$ across policies (see Section 5.4), which represents the effectiveness of reading the first $n$ comments given the policies being evaluated. These values can be calculated for a single debate, or across an ensemble of debates, for a given policy. For the latter case, we can compare summary statistics of the distributions of values of the averages or $J_n$'s and use statistical tests to see which policy has a greater such summary statistic.

As mentioned earlier, if for simplicity one does not wish to incorporate argument mining techniques, then these calculations can only make sense if we do have a dataset of "clean" online debates, i.e. where each comment qualifies as a self-contained argument, and each reply is classified as either an attack or a support. Fortunately, the online debating platform \textit{Kialo} \footnote{https://www.kialo.com/, last accessed 26/8/2020.} provides such a "clean" dataset, thanks to its moderation policy that ensures every comment is a concise and relevant argument to the topic being discussed that is replying either positively (supporting) or negatively (attacking) to another comment (or no such reply if it is the root comment of a discussion). We have mined Kialo data until March 2019 (see Section 4), and processed this data via our proposed pipeline. Subject to suitable aggregation and identification of what "likes" and "time" mean for each comment (see Section 5.2), the pipeline shows that reading Kialo debates from the most to least "liked" comment displays more actually winning arguments, \textit{on average}, than reading Kialo debates from the earliest comment to the most recent comment. We close by conjecturing why this is and speculate on how robust this result is to less "clean" online debates.

1.3. Structure of this paper

In Section 2, we will summarise the relevant background in argumentation theory and related work. In Section 3, we will explain the pipeline and clarify its assumptions and limitations. In Section 4, we overview Kialo, explain how it validates the pipeline’s assumptions, explain how we have mined the data, and offer summary statistics. In Section 5, we articulate the comment sorting policies based on likes and time for Kialo, and have the pipeline measure which policy is more effective. Section 6 concludes with future work.
2. Technical background and related work

In this section we recap relevant aspects of abstract argumentation theory and bipolar argumentation theory. We then give a brief survey of argument mining and its application to online debates.

2.1. Abstract argumentation theory

From [15]: An abstract argumentation framework (AF) is a digraph \( \langle A, R \rangle \), where \( A \) is the set of arguments under consideration and \( R \subseteq A^2 \) is the defeat relation, where \( A^2 := A \times A \) (and similarly, \( X^n \) denotes \( X \times \cdots \times X \) \( n \) times). We write \( R(a, b) \) to denote \( (a, b) \in R \), which means \( a \) defeats (disagrees with) \( b \). In what follows let \( a, b \in A \) and \( S \subseteq A \). Let \( S^+ := \{ b \in A \mid (\exists a \in S) R(a, b) \} \), \( S^- := \{ b \in A \mid (\exists a \in S) R(b, a) \} \) and \( a^\pm := \{ a \}^\pm \). Define the neutrality function, which maps sets of arguments to sets of arguments, to be \( nf : P(A) \mapsto P(A) : S \mapsto A - S^+ \), i.e. the set of arguments not defeated by \( S \). We say \( S \) is conflict-free (cf) iff \( S \subseteq nf(S) \), and \( S \) is self-defending (sd) iff \( S \subseteq nf ons(S) \), where \( o \) denotes composition of functions (if well-defined).

We say \( S \) is admissible iff \( S \) is both cf and sd. There are four basic notions of “winning arguments”, each of which builds on admissibility [15]. We say \( S \) is a complete extension iff \( S \) is admissible and that \( nf \circ nf(S) \subseteq S \). We say \( S \) is a preferred extension iff \( S \) is a \( \subseteq \)-maximal complete extension. We say \( S \) is a stable extension iff \( nf(S) = S \). We say \( S \) is the grounded extension, denoted \( G \), iff it is the \( \subseteq \)-least complete extension; \( G \) always exists, is unique, and is empty iff the set of unattacked arguments (i.e. \( a \in A \) where \( a^- = \emptyset \)) is empty [37, Corollary 6.10].

Recall that if an AF is well-founded, i.e. that there exists no sequence of arguments \( \{a_i\}_{i \in \mathbb{N}} \) such that (\( \forall i \in \mathbb{N} \)) \( R(a_{i+1}, a_i) \) [15, Definition 29], then \( G \) is the unique stable, preferred and complete extension, i.e. there is only one set of winning arguments [15, Theorem 30]; this is certainly the case when \( \langle A, R \rangle \) is a tree [37, Corollary 2.46].

Example 2.1. Let \( A = \{a, b, c, d, e\} \) and \( R = \{(c, b), (c, a), (d, b)\} \). \( \langle A, R \rangle \) is well-founded with \( G = \{c, d, e\} \). This AF is shown in Fig. 1.

2.2. Bipolar argumentation theory

Bipolar argumentation theory extends AFs with a support relation [9]. A bipolar argumentation framework (BAF) is a structure \( \langle A, R_{att}, R_{sup} \rangle \) where \( A \) is the set of arguments, \( R_{att} \subseteq A^2 \) denotes when two arguments attack each other, and \( R_{sup} \subseteq A^2 \) denotes when two arguments support each other, and \( R_{att} \cap R_{sup} = \emptyset \). A BAF can be visualised as a digraph where edges are either red (\( R_{att} \)) or green (\( R_{sup} \)).

There are principled ways to absorb \( R_{sup} \) into \( R_{att} \), resulting in an associated AF \( \langle A, R \rangle \). Following [9], for arguments \( a, b \in A \), we say that a support path from \( a \) to \( b \) is a path in the underlying digraph that only traverses support edges. We write \( a \to_{sup} b \) iff a support path exists from \( a \) to \( b \). We say a support-defeats \( b \) iff (\( \exists c \in A \)[(\( a \to_{sup} c \), \( R_{att}(c, b) \))] and a indirectly defeats \( b \) iff (\( \exists c \in A \)[\( R_{att}(a, c) \), \( c \to_{sup} b \)]. The associated AF of a given BAF is \( \langle A, R \rangle \), where \( R(a, b) \) iff \( a \) support-defeats \( b \) or \( a \) indirectly defeats \( b \). The winning sets of arguments for a BAF are the extensions of its associated AF.
Example 2.2 (Example 2.1 continued). Figure 2 illustrates the BAF where $A = \{a, b, c, d, e\}$, $R_{\text{att}} = \{(c, b)\}$ and $R_{\text{sup}} = \{(d, c), (e, b), (b, a)\}$.

The support paths of length 1 in this BAF are $(e, b)$, $(b, a)$, and $(d, c)$. The support paths of length 2 in this BAF consist of only $(e, b, a)$. Therefore, $e \rightarrow_{\text{sup}} b$, $b \rightarrow_{\text{sup}} a$, $d \rightarrow_{\text{sup}} c$ and $e \rightarrow_{\text{sup}} a$. As $c$ attacks $b$, we can see that $c$ indirectly defeats $a$, and $d$ support-defeats $b$. Therefore, the associated AF of this BAF has the same arguments, but the defeat relation is $R = \{(c, b), (c, a), (d, b)\}$. This resulting AF is already illustrated in Fig. 1. The set of winning arguments for this BAF is also $G = \{c, d, e\}$.

2.3. Argument mining and the analysis of online debates

Argument mining (e.g. [23,24]) seeks to identify and extract structured, self-contained arguments and inter-argument relationships from natural language text to (amongst many other applications) resolve conflicting points of view expressed in online debates (e.g. [6,7]). However, in order to successfully apply argumentation theory to real data, one would need a relatively “clean” dataset of online debates. This is non-trivial, because most online debates are very “noisy”. For example, Twitter Tweets are often too short to be treated as self-contained arguments so a preprocessing step is required to determine whether Tweets can be treated as such [3]. Further, one must formally represent the relationships between the mined arguments. In the context of applying BAFs to model online debates, we are interested in classifying whether a reply from comment $a$ to comment $b$ (assuming both $a$ and $b$ can be treated as arguments) is attacking or supporting. Existing techniques include textual entailment, which infers relevance and contradiction between pieces of natural language (e.g. [2,12]) and hence identify attacks between comments. Another way of identifying attack or support can be to calculate affective or emotional scores from key words in the comments using Linguistic Inquiry and Word Count [29,34], such that replies between comments with very similar such scores can be related to supports. A more sophisticated method for identifying attacks and supports involves deep learning [11]. Therefore, there are many argument mining techniques at our disposal if one wishes to investigate a “noisy” online debate dataset, where it cannot be assumed that each comment made by users constitutes a well-defined argument, or that every reply is either an attack or a support. We envisage that such techniques are used to first render the “noisy” online debate “clean”, then one can use the usual methods in identifying the winning set(s) of arguments.

In the next section, we will apply ideas from argumentation theory to articulate a pipeline that allows one to rank the effectiveness of various comment sorting policies in online debates, assuming that the input data is already “clean”, either because it is “naturally clean”, or that it has been “cleaned up” by techniques such as those just mentioned. Therefore, the pipeline does not preclude the use of such techniques. We will outline how such techniques can be integrated into this pipeline in Section 6.
3. **A pipeline to evaluate comment sorting policies**

In this section, we articulate the pipeline, based in argumentation theory, that can evaluate how best to sort comments for readers who may not have read the entire debate.

3.1. **The stages of the pipeline**

The pipeline consists of the following stages, where the first two stages involve validating our assumptions on the data.

1. We input to the pipeline an online debate. The nodes are comments made in the debate that possess some *attributes*, e.g. upvotes, downvotes, Retweets (in the case of Twitter), time of posting, the author’s username... etc. The edges are replies, such that for comments *a* and *b*, the edge points from *b* to *a* iff *b* replies to *a*. As the online debates we will focus on in this work are trees, we will restrict to the case of trees as inputs to the pipeline, i.e. that each comment beyond the first can only reply to one other comment, and the digraph is weakly connected. Further, we do not consider time-evolution of debates, i.e. the debate accepted by the pipeline is a snapshot of the comments at that given point in time.

2. We make the assumption that the digraph is “clean”, in that every comment can be seen as a self-contained argument, and every reply is either an attack or a support. Therefore, the debate digraph is a BAF $⟨A, R_{att}, R_{sup}⟩$ that is also a tree. Let $A$ contain $N ∈ \mathbb{N}^+$ arguments.

3. We calculate the grounded extension, $G ⊆ A$, for this BAF, via its associated AF (this AF may not be connected, see Examples 2.1 and 2.2). We interpret this as the set of actually winning arguments at that point in time of the debate. The choice to calculate $G$ (as opposed to, e.g. the preferred extensions) is justified in that if the underlying debate digraph is well-founded, then its BAF and hence AF are also well-founded [9], which has only one extension [15, Theorem 30] – the grounded extension. Further, calculating $G$ takes polynomial time [16]. Despite these desirable formal properties, the pipeline thus assumes that human readers are “sufficiently rational” such that they would agree that arguments in the grounded extension are indeed winning.

4. We now explain what we mean by *comment sorting policy with respect to a given (numerical) attribute*, and provide some simple examples; examples based on real data will be motivated and defined in Section 5.2. Our input BAF is a tree of debate comments and replies. There are many ways to read a debate. The pipeline currently considers the following way: start with the root comment, and at a given comment, read the unread reply with the most or least of that numerical attribute, until a comment with no replies has been reached (a *leaf* node), and then backtrack to the most recent comment where there is more than one reply, and continue. This is depth-first search (DFS), and sorts the comments in $A$ into a list $(a_1, a_2, \ldots, a_N)$, where $a_1$ is the root argument that starts the debate and does not reply to anything.

Why is this method of reading comments plausible? We assume that a reader who is interested in the debate would start with the opening argument and read sequentially along the replies without missing any intermediate arguments, in order to fully understand each thread of the debate. Upon reaching an argument with multiple replies, the online debate platform will rank the subsequent debate replies by the chosen attribute (e.g. number of likes), therefore the reader will start reading the debate along the path with the greatest or least values of this attribute. Indeed, this is enforced
to some extent in real online discussion platforms such as comments to articles in Disqus\(^4\) or The Daily Mail (a UK newspaper).\(^5\) We now give two examples:

**Example 3.1** (Example 2.2 continued). This example sorts the arguments to represent reading from the earliest to the most recent following DFS. Consider the BAF from Example 2.2, where \(A = \{a, b, c, d, e\}\), \(R_{\text{att}} = \{(c, b)\}\) and \(R_{\sup} = \{(d, c), (e, b), (b, a)\}\). The root argument – the argument that starts the debate and hence does not reply to anything else – is \(a\), and the leaf arguments – those that have not received any replies at this point in time – are \(d\) and \(e\). Suppose that each argument \(x \in A\) has a timestamp \(t_x \in \mathbb{R}\) (a real number), such that the larger \(t_x\) is, the more recent the argument has been posted to the debate. We assume that if argument \(y\) replies to argument \(x\), then \(t_y > t_x\).

A reader would start at the root, \(a\), which has received only one reply, \(b\) and \(e\) in turn receive two replies, \(c\) and \(e\). Assume that \(e\) has been posted prior to \(c\) (i.e. \(t_e < t_c\)), then the next argument after \(e\) is \(b\), and we reach a leaf argument. We then backtrack to \(b\), where the remaining unread reply is \(c\), from which one reads \(d\). This covers all five arguments, and the order of reading through the comment sorting policy from earliest to most recent following DFS is \((a, b, e, c, d)\).

From now on, we make the assumption of time coherence of a debate network – that if argument \(b\) replies to argument \(a\), then \(t_b > t_a\), i.e. a comment must exist first before it can be replied to. Furthermore, we will use the terms root and leaf as defined in Example 3.1.

**Example 3.2** (Example 3.1 continued). This example sorts the arguments to represent reading from the most liked to the least liked following DFS. Consider again the BAF in Example 2.2. Suppose further that argument \(c\) is more liked than argument \(e\). The resulting sorting policy based on DFS and reading from most to least liked will output \((a, b, c, d, e)\), which is distinct from \((a, b, e, c, d)\) in Example 3.1.

(5) As shown in Examples 3.1 and 3.2, comment sorting policies order the arguments in \(A\) into a list \((a_1, a_2, \ldots, a_N)\). Let \(n \in \mathbb{N}\) denote the number of comments read by a reader under this policy. Clearly, \(0 \leq n \leq N\), so it is possible for the reader to read nothing \((n = 0)\) or everything \((n = N)\). For each such \(n\), we consider the first \(n\) arguments of the list \((a_1, a_2, \ldots, a_N)\), and obtain the induced sub-BAF. Each induced sub-BAF has its own grounded extension \(G_n\), each of which represents the seemingly or provisionally winning arguments given that a reader has only read \(n\) comments from a policy. Clearly, \(G_0 = \emptyset\) (the empty BAF has the empty set of winning arguments), and \(G_N = G\).

(6) We quantitatively compare \(G\) and \(G_n\) using the Jaccard coefficient: for finite sets \(A\) and \(B\),

\[
J(A, B) := \frac{|A \cap B|}{|A \cup B|}
\]

\(\in [0, 1]\), where for any set \(X, |X|\) is the number of elements in \(X\), (3.1)

and \(J(\emptyset, \emptyset) := 0\); this is the normalised overlap between two sets. For each \(n\), we calculate \(J_n := J(G, G_n)\), the proportion of actually winning arguments present in \(G_n\). A large \(J_n\) means that for

\(^4\)E.g. see \url{https://www.quantamagazine.org/physicists-debate-hawkings-idea-that-the-universe-had-no-beginning-20190606}, last accessed 26/8/2020. Disqus has appropriately indented comments in the user-interface (UI) to denote the depth of each comment within the debate tree, and sorting the comments from e.g. most to least liked follows this DFS display pattern.

\(^5\)E.g. see \url{https://www.dailymail.co.uk/news/article-6894539/Meghan-Markle-snubs-Queens-doctors-birth-doesnt-want-men-suits.html}, last accessed 26/8/2020. In contrast, the comments, although sortable by DFS, only allows replies of depth 2 (i.e. replies to the news article, and replies to replies), and displays only the first two comments amongst the depth-2 replies by default.
Table 1
The results of applying steps (5) and (6) of the pipeline in Example 3.3

<table>
<thead>
<tr>
<th>n</th>
<th>Comments Read</th>
<th>Actual winners $G$</th>
<th>Apparent winners so far, $G_n$</th>
<th>Jaccard coefficient, $J(G, G_n)$</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>()</td>
<td>[c, d, e]</td>
<td>$\emptyset$</td>
<td>0</td>
</tr>
<tr>
<td>1</td>
<td>(a)</td>
<td>[c, d, e]</td>
<td>[a]</td>
<td>0</td>
</tr>
<tr>
<td>2</td>
<td>(a, b)</td>
<td>[c, d, e]</td>
<td>[a, b]</td>
<td>0</td>
</tr>
<tr>
<td>3</td>
<td>(a, b, c)</td>
<td>[c, d, e]</td>
<td>[a, b, c]</td>
<td>0.2</td>
</tr>
<tr>
<td>4</td>
<td>(a, b, c, d)</td>
<td>[c, d, e]</td>
<td>[e, c]</td>
<td>0.666 ...</td>
</tr>
<tr>
<td>5</td>
<td>(a, b, c, d, e)</td>
<td>[c, d, e]</td>
<td>[e, c, d]</td>
<td>1</td>
</tr>
</tbody>
</table>

Example 3.3 (Example 3.1 continued). Reading the BAF from Example 2.2 from the earliest to the latest comment following DFS, we obtain the reading order $(a, b, e, c, d)$. The reader may read $0 \leq n \leq 5$ of these arguments. Table 1 gives the results of Steps (5) and (6) of the pipeline for this case.

Example 3.4 (Example 3.2 continued). Reading the BAF from Example 2.2 from the most liked to the least liked comment following DFS, we obtain the reading order $(a, b, c, d, e)$. We calculate the same table (see Table 2) in Example 3.3 for this case.

We illustrate the above six steps, for a fixed value of $n$, with Fig. 3 [38]:

3.2. Formal properties of the pipeline

The following are straightforward consequences of our pipeline.
Corollary 3.5. In our pipeline, illustrated in Fig. 3, $J_0 = 0$ for all policies.

Proof. If $n = 0$, then the induced BAF is the empty BAF, which is associated with the empty AF, which has $G_0 = \emptyset$. Equation (3.1) implies $J(G, \emptyset) = 0$. □

Note that the converse to Corollary 3.5 is false, as indicated in Example 3.3 for $n = 2$, say. This means it is possible to read a non-zero portion of the debate and think (at least normatively) that the actually winning arguments are all losing, therefore “totally missing the point”.

Corollary 3.6. Let $N$ be the number of arguments of our input BAF, then $J_N = 1$ for all policies.

Proof. The induced BAF is the original BAF itself, hence $G_N = G$ and $J(G, G) = 1$. □

Intuitively, Corollary 3.5 states that if you read nothing you get nothing, and Corollary 3.6 states that if you read everything you get everything, regardless of the comment sorting policy. One further result concerns the case where all edges are supports and the policy is based on DFS.

Theorem 3.7. Consider an input BAF $\langle A, R_{att}, R_{sup} \rangle$ of $N$ nodes and a comment sorting policy based on DFS. We have that $R_{att} = \emptyset$ iff $(\forall 0 \leq n \leq N) J_n = \frac{n}{N}$.

Proof. $(\Rightarrow)$ If $R_{att} = \emptyset$, then the associated AF $\langle A, R \rangle$ of this BAF (Section 2.2) will satisfy $R = \emptyset$. Therefore, for each $0 \leq n \leq N$, $G_n = \{a_1, a_2, \ldots, a_n\}$. As $G = \{a_1, a_2, \ldots, a_N\} \supseteq G_n$, we have $G \cap G_n = G_n$ and $G \cup G_n = G$. It follows from Equation (3.1) that $J(G, G_n) = \frac{n}{N}$.

$(\Leftarrow)$ (contrapositive) Assume that there is some attacking edge $R_{att}(a, b)$. WLOG we choose an attacking edge where $a \in G$; this is possible because there is a unique path between any two nodes in a tree, so we choose the attacking edge where either $a$ is a leaf (and hence unrebutted), or there is a support path (Section 2.2) from a leaf to $a$. Therefore, $b \notin G$. But as we are reading DFS from the root node outwards (Section 3.1), as $n$ increases, we will end up reading the losing argument $b$ before the actually winning argument $a$. Therefore, $(\exists 0 \leq n \leq N) J(G, G_n) < \frac{n}{N}$, as $G_n$ will have an argument not in $G$. □

This means that in a debate where there is no disagreement, the more one reads, the more of actually winning arguments is obtained, because every additional argument read is actually winning. Further, this strictly monotonically increasing situation is only possible in the case where every reply is supporting.

3.3. How are comment sorting policies compared?

As discussed so far, the pipeline takes as input a BAF and a comment sorting policy, and returns a vector $J := (J_0 = 0, J_1, J_2, \ldots, J_{N-1}, J_N = 1)$, where each $J_n \in [0, 1]$. The larger the $J_n$, the larger the proportion of actually winning arguments one is exposed to by only reading the first $n$ comments according to that comment sorting policy. How can we aggregate the components of $J$ for each policy in order to compare different policies? One idea is to take the average across all $N + 1$ components:

Definition 3.8. Given a policy $P$ that has a $(N + 1)$-component vector of Jaccard coefficients $J := (0, J_1, J_2, \ldots, J_{N-1}, J_{N-1}, 1)$, its average Jaccard is $J := \frac{1}{N+1} \sum_{n=1}^{N} J_n$.

Intuitively, as the Jaccard coefficient is a measure of how many actually winning arguments the apparently winning arguments contain, having read the debate up to a certain point, $J$ for a policy is the typical proportion of actually winning arguments the policy exposes. We now give an example.
Example 3.9 (Examples 3.3 and 3.4 continued). For the BAF in Example 2.2, sorting from the earliest comment to the most recent comment gives an average Jaccard of $\langle J \rangle_{\text{time ascending}} = \frac{13}{45} \approx 0.31$. Sorting from most to least liked gives an average Jaccard of $\langle J \rangle_{\text{likes descending}} = \frac{1}{3} \approx 0.33$. Therefore, for this BAF, $\langle J \rangle_{\text{likes descending}} > \langle J \rangle_{\text{time ascending}}$. It is more effective to sort the comments from most to least liked, compared to sorting from earliest to latest.

The following corollary of Theorem 3.7 shows for the special case of debates where there is no disagreement, every policy based on DFS gives an average Jaccard of 0.5. This makes sense as we are taking the centre of mass of the diagonal line that joins $(\frac{n}{N}, J_n) = (0, 0)$ and $(1, 1)$.

Corollary 3.10. If the input BAF satisfies $R_{\text{alt}} = \emptyset$, then all policies based on DFS satisfy $\langle J \rangle = 0.5$.

Proof. By Theorem 3.7, for any sorting policy based on DFS, $(\forall 0 \leq n \leq N) J_n = \frac{n}{N}$. Therefore,

$$\langle J \rangle = \frac{1}{N+1} \sum_{n=1}^{N} \frac{n}{N} = \frac{1}{N(N+1)} \sum_{n=1}^{N} n = \frac{1}{N(N+1)} \frac{N(N+1)}{2} = \frac{1}{2},$$

where we have used a basic result concerning the sum of the first $N$ non-zero natural numbers. □

3.4. Summary of the pipeline

The first contribution of this paper is a pipeline that accepts as input an online debate represented as a BAF, and a comment sorting policy, and calculates how this policy, on average, exposes the reader to the actually winning arguments, especially when the reader has not read the entire debate. We have given a running example of how the pipeline works (Examples 2.1 to 3.9).

However, the fundamental assumption of the pipeline – that a real-life online debate can be represented as a BAF – is suspect. One cannot just assume that comments posted on online debates are self-contained arguments, nor that every reply can be classified as an attack or a support. Of course, the pipeline does not preclude the use of argument mining techniques (Section 2.3) to structure online debates into a BAF for our pipeline to ingest (Section 6). In the next section, we present a dataset that we have mined, which seems to satisfy the property of being “clean”, and thus being immediately suitable for input into the pipeline as a proof of concept of how might various comment sorting policies be evaluated.

4. Kialo

We now overview Kialo, an online debating platform. We argue that its moderation policy ensures that the debates it hosts are sufficiently “clean” to be inputted into our pipeline for analysis. We summarise how the debates were mined and cleaned into BAFs for our dataset.

4.1. What is Kialo and why is Kialo “clean”?

Kialo is an online debating platform that helps people “engage in thoughtful discussion, understand different points of view, and help with collaborative decision-making.”\(^6\) In a Kialo debate, users submit

claims. The starting claim is a thesis, which has at least one, but possibly more, indexing tags. Further claims then reply to existing claims, and are classified as either pro or con the claim being replied to.

**Example 4.1.** Consider the Kialo discussion with thesis, “Internet companies are wrong in denying services to white supremacists.” One example pro argument for the thesis is, “Denying service to white supremacists might lead to extensive online censorship of non-mainstream views.” One example con argument is, “It is in the interest of Internet companies to deny services to white supremacists.”

Kialo debates are “clean” and easily-represented as BAFs without the use of argument mining tools (Section 2.3). Firstly, Kialo has a strict moderation policy, which aims to keep users engaged and promotes a well-structured discussion.\(^8\) Central to this is the enforcement of writing good claims. For example, each claim should be concise, self-contained, based on logic and facts, and make a single point that is relevant to the debate topic – in other words, claims are arguments.\(^9\) Claims should not be questions, or mere comments, or duplicated within a discussion – this excludes the possibility that claims are insults, say. Discussions are structured such that more general claims are closer to the root, and more specific claims are closer to the leaves. The conciseness of each claim allows for fine-grained support or attacks when replied to.\(^10\) Claims that fail to meet such criteria are marked for review.\(^11\) Either the claim is deemed unsuitable and therefore removed from the debate, or that the claim is conditionally approved after some discussion between the claim’s author and the moderators, and that the claim’s author must make satisfactory improvements for the claim to be published.\(^12\) This validates the pipeline’s assumption that all claims made in Kialo debates can be treated as an argument.

In Kialo, every claim \(b\) that is replying to another claim \(a\) is classified by the author as pro or con \(a\); the polarity of the reply is also suitably moderated as part of auditing the claim. Therefore, all replies are classified as either attacking or supporting. This validates our assumption that Kialo debates can be modelled as BAFs. Further, the mechanism of how claims reply to at most one other claim and that each claim is moderated to ensure that it makes a unique point mean we can assume Kialo debates are trees.

In summary, Kialo is an online debating platform. Its highly moderated nature validates our assumptions that every claim made in Kialo debates can be treated as self-contained arguments. Further, every reply between two claims in a Kialo debate must either be a support (pro) or an attack (con). Lastly, the structure of Kialo debates guarantees that the underlying digraph is a tree. It is in this sense that Kialo debates are sufficiently “clean” for us to represent them straightforwardly as BAFs without the need for further data cleaning methods.

### 4.2. Mining Kialo

We scraped all Kialo debates dated up until March 2019 as follows: we use the publicly visible Kialo application programming interface (API), which is used by Kialo’s front end, to acquire all the available tags used to tag individual debates; this resulted in a list of over one thousand unique tags. We then used

---

12 Refer to the URL in Footnote 8.
the same API to perform a tag-based search for debates using the previously acquired list of tags. For each debate we obtained its claims, replies, and various claim attributes such as the username of the author, its text, time of posting, and votes (related to likes, see Section 5.2). This returned 1,056 debates. We then use independent means of verifying this dataset [4], which gives us a high degree of confidence that this is almost all of the debate activity on Kialo, as of March 2019.

4.3. The resulting dataset

When the debates dataset was cleaned, we noticed that some nodes (less than 1% of the total) have empty text and therefore could not be considered arguments; this is possibly due to how the data is stored in the back end of Kialo has failed to synchronise with moderators asking for a claim to be removed. Further, some reply edges violate time coherence (Section 3.1). We delete both such nodes and edges and focus on the resulting sub-discussions, on the assumption that Kialo’s moderation policy implies that sub-discussions are also self-contained, despite being separated from the thesis, due to the conciseness of each reply. This results in 49,040 weakly-connected BAFs that are also trees, containing 1 to more than 500 arguments. By excluding BAFs with 9 arguments or less, we obtain 4,365 BAFs to analyse with our pipeline. These BAFs have an average of 28.71 arguments (sample standard deviation (SSD) 35.21). Amongst these BAFs, the in-degrees of all nodes present range from 0 (leaves) to 58, with a mean in-degree of 0.965 (SSD 2.00). Amongst all root nodes, the average in-degree is 7.34 (SSD 4.58). This gives some idea of the typical activity of discussion in the average debate.

Example 4.2 (Example 4.1 continued). We can visualise the debate in Fig. 4. This has 115 arguments where 57% of replies are attacks, and 43% of replies are supports.

![Fig. 4. The debate as a BAF from Example 4.2, where red edges are attacks and green edges are supports.](image)

13The anonymised BAF graphs from the Kialo data would be shared upon reasonable request for research and reproducibility.
5. Analysing Kialo debates with our pipeline

We now input our Kialo dataset (Section 4) into our pipeline (Section 3) and evaluate four policies.

5.1. Validating the pipeline assumptions

As stated in Section 4.1, Kialo debates are sufficiently clean such that they are readily representable as BAFs due to its design and moderation policy. Further, in our dataset, all debates are trees hence all debates (and all induced subgraphs thereof) have a well-defined set of winning arguments via the grounded extension that the pipeline calculates. This verifies Steps (1) to (3) of the pipeline (Section 3.1), apart from defining the policies to sort comments by. Lastly, we do not need to perform further natural language processing, because the BAFs are already given.

5.2. Which policies to evaluate?

To apply Step (4) of the pipeline to Kialo data, we still assume that an interested reader would read each replying comment from the root in a DFS manner as this models how the reader is getting whole conversation threads without skipping comments before backtracking to the next possible thread.

For the purposes of illustrating how the pipeline works on “cleaned” real data in this paper, we compare sorting along two attributes: likes and time; other attributes and policies will be considered in future work (Section 6), as these two attributes are readily available from Kialo data.

Time refers to the time of posting the claim, while likes is a measure of how “good” the claim is. In our dataset, each comment in a Kialo debate is already time-stamped, hence time for each comment is well-defined. However, likes is not well-defined. Conceptually, the closest attribute available in Kialo comments is the number of votes, which is a five-tuple $v = (v_0, v_1, v_2, v_3, v_4)$, where each $v_k \in \mathbb{N}$ is the number of votes made by users for the $k$th impact category of the comment. Intuitively, impact is a function of the truth, persuasiveness and relevance of the claim to the parent it is replying to. Claims are assigned votes close to the end of the discussion (as deemed by the moderator), and when each user votes, the other votes claims have already received are hidden, which avoids bias. Impact category 0 means the comment has no impact, while category 4 denotes very high impact.

How can we relate the five-vector of votes $v$ to a single number to sort claims by, which is also a proxy for likes? We use the following weighted average: define a weight vector $w := (w_0, w_1, w_2, w_3, w_4) \in (\mathbb{R}_0^+)^5$, which weighs each impact category by the corresponding non-negative weight $w_k$. Given such a $w$, define $\mu(w) := \frac{1}{5} \sum_{k=0}^{4} w_k$; if there is no risk of ambiguity about $w$ we will write $\mu$ instead. Let $\cdot$ denote the usual (Euclidean) dot product. We define a normalised weighted average, parameterised by $w$, as

$$nwa_w(v) := \frac{v \cdot w}{\sum_{k=0}^{4} v_k}$$

if $v \neq 0$, and $nwa_w(0) := \mu(w), \quad (5.1)$

where 0 is the null vector in $\mathbb{R}^5$. This is a well-defined function from $\mathbb{N}^5$ to $\mathbb{R}$.

Example 5.1. If a claim has $v = (1, 0, 0, 4, 3)$ and $w = (1, 1, 1, 1, 10)$, then $nwa_w(v) = \frac{35}{2} = 4.375$.

\[14\text{See https://support.kialo.com/hc/en-us/articles/115001605352-Using-Voting, last accessed 26/8/2020.}\n
Notice if \( w = 0 \), then \( (\forall v \in \mathbb{N}^5) \ nwa_w(v) = 0 \). If all components of \( w \) are equal to 1, then \( (\forall v \in \mathbb{N}^5) \ nwa_w(v) = 1 \). Further, if all \( w_k \) of \( w \) satisfy \( w_k \leq 1 \), then \( (\forall v \in \mathbb{N}^5) \ nwa_w(v) \leq 1 \). Therefore, we consider \( w \) where all components \( w_k \) lie strictly between 0 and 1.

We now define the following two \( w \)'s to form two “sort by likes” policies. Informally, the first \( w \) is \textit{a priori} and the second is data-driven. Suppose we interpret these five impact categories in terms of a rational reader’s degrees of belief about the truth of the claim (e.g. [21]), so \( v_0 \) denotes the number of votes by users who think the claim is false... etc. and \( v_4 \) denotes the number of votes by users who think the claim is true. In this case, we let \( w_e := (0, 0.25, 0.5, 0.75, 1) \), which weighs the five categories equally in the sense that moving up between two categories corresponds to an equal amount of increase in the voters’ degrees of belief concerning the claim’s truth. We call \( nwa_w(v) \) the \textit{rational likes} of \( v \).

**Example 5.2.** If a claim has \( v = (1, 0, 0, 4, 3) \), then its rational likes is 0.75, by Equation (5.1).

One criticism of rational likes is that this idea of equal weight increases across the five impact categories is unjustified in that it assumes all readers interpret these impact categories rationally. Another approach is to weigh each impact category according to the number of votes each has received across all categories is unjustified in that it assumes all readers interpret these impact categories rationally. We call \( nwa_w(v) \) the \textit{empirical likes} of \( v \).

**Example 5.3.** If a claim has \( v = (1, 0, 0, 4, 3) \), then its empirical likes is 0.2279, by Equation (5.1).

Therefore, we enhance all nodes in each BAF of our dataset with rational and empirical likes, based on their votes vector. We define the following four policies to evaluate with our pipeline (Section 3.1):

1. **Policy 1:** Non-DFS Ascending Time sorts the comments from oldest (smallest Unix time) to most recent (largest Unix time). This is “Non-DFS” because we do not follow DFS as described in Section 3.1 to serve as a baseline for the next policy; not following DFS means that a reader may move between discussion threads rather than finishing a complete thread first.

2. **Policy 2:** Ascending time with DFS always choose the lowest value of (Unix) time that has not yet been chosen in DFS when there is a choice of replies – this is well-defined as there are no ties in the timestamp of comments. This displays the comments from the earliest posted to the most recently posted while preserving each branch of the debate.

3. **Policy 3:** Descending rational likes with DFS, i.e. always choose the highest value of rational likes that has not yet been chosen in DFS. This sorts the comments from most to least liked following the debate branch. If there are ties, we break the tie with respect to ascending time, i.e. when choosing between two equally rationally liked comments, choose the earlier one first.

4. **Policy 4:** Descending empirical likes with DFS is analogous to Policy 3, but replaces rational likes with empirical likes, with the same tie-breaking mechanism.

5.3. Results

We now measure the effectiveness of the above four policies as described in Steps (5) and (6) of Section 3.1. For each debate in our dataset, we calculate and plot \( \{ (\frac{n}{N}, J_n) \}_{n=1}^{N} \) for each policy. Plotting \( J_n \) against \( \frac{n}{N} \) instead of \( n \) will allow us to make comparisons across debates of different sizes, i.e. in terms of the \textit{fraction} of the debate read with respect to a policy. The following example shows each policy’s \( \{ (\frac{n}{N}, J_n) \}_{n=1}^{N} \) plot for the debate in Example 5.4.
Example 5.4 (Example 4.2 continued). The data cleaning process has assigned the index 2,768 to this debate. As \( n \) increases from 0 to \( N \), Fig. 5 displays how \( J_n \) changes with respect to \( n/N \).

Visually, it is difficult to judge how each policy differs as all cluster around the diagonal. For each debate, we plot the detrended Jaccard, \( DJ_n := J_n - \frac{n}{N} \) against \( \frac{n}{N} \), for \( 0 \leq n \leq N \). Notice \( DJ_0 = DJ_N = 0 \).

Example 5.5 (Example 5.4 continued). Figure 6 plots the detrended Jaccard vs. \( \frac{n}{N} \) for Example 4.1.

Visually, all policies are the same when reading up to \( \sim 2\% \) of this debate. Further, Policy 3 (descend-
ing rational likes with DFS) and Policy 2 (ascending time with DFS) seem to track each other, and both seeming to beat Policy 3 (ascending time without DFS), while Policy 1 (ascending time, no DFS) seems to capture the least proportion of the actually winning arguments.

The conclusions we have drawn so far only apply to debate 2,768 (Examples 4.2 to 5.5). We are now interested in using our entire dataset of 4,365 BAFs to determine which of our four policies – Policies 1 and 2 are based on time (with or without DFS) and Policies 3 and 4 are based on likes (rational and empirical, with DFS) – is the most effective, i.e. maximises $J_n$. We will take the average $DJ_n$ across $0 \leq n \leq N$ for each debate under each policy, as it is easy to show that this is equal to $\langle J \rangle - \frac{1}{2}$ (the proof is analogous to that of Corollary 3.10). Therefore, comparing average values of $DJ_n$ is the same as comparing $\langle J \rangle$. We then plot the distribution of such values for each policy.

5.4. For Kialo, sorting by likes is, on average, better than sorting by time

We repeat the calculation in Example 5.5 for all 4,365 BAFs in our dataset. Each version of Fig. 6 for each BAF will have some pattern of fluctuations about the horizontal for each policy. Instead of describing these fluctuations directly, we evaluate the policies over our dataset of debates in two ways: (1) across the entire range of $\frac{n}{N}$ (Section 3.3), and (2) evaluating each policy by calculating $J_5$, which models the effect of reading the top five claims given the policy. Admittedly five is arbitrary; we assume that a hurried reader only has time to read five claims. The distributions of the averaged $DJ_n$ across all four policies for (1) are shown in Fig. 7. The analogous distributions for (2) are shown in Fig. 8.

Visually, we can see that Policy 1 is the worst policy to sort by, according to its averaged $DJ_n$ distribution. Policy 2 seems much better, while Policies 3 and 4 (both based on likes), overlap a lot. We can make this more precise with the summary statistics of the distributions in Fig. 7 in Table 3.

![Histogram](image)

Fig. 7. The $DJ_n$ distributions for all policies averaged over $\frac{n}{N}$ (“one segment”), for all 4,365 trees.
We can see that comparing the medians and means of the four policies does support that sorting by likes is better than sorting by time, and further, that time with DFS is better than time without DFS. Let us make this even more precise with an appropriate statistical test. Figure 7 displays skewed distributions, therefore we should make no assumptions on the underlying population distribution for average $DJ_n$ per policy (non-parametrics). Indeed, performing a Kolmogorov–Smirnov test [22,33] allows us to conclude that none of these distributions are Gaussian, and hence we should compare medians rather than means.

As this is a non-parametric situation, we apply the Mann–Whitney test [25] to compare whether the observed distributions of average $DJ_n$ are sampled from the same underlying population distribution (our null hypothesis), or different such distributions (our alternative hypothesis). At a significance level of $\alpha = 0.05$, we perform $6 = \binom{4}{2}$ such tests. The results for the data plotted in Fig. 7 are shown in Table 4.

For Fig. 7, Policy 2 beats Policy 1 (average $DJ_n = -0.067 > -0.107$), Policies 3 and 4 beat Policy 2 (average $DJ_n = -0.052, -0.053 > -0.067$), while Policies 3 and 4 are not significantly better than each other (average $DJ_n = -0.052 \sim -0.053$). Overall, sorting by likes, averaged over all debates in our dataset, beats sorting by time.

We now perform the same calculation for the case where instead of averaging $DJ_n$ over all of $0 \leq n \leq N$, we evaluate each policy using $J_5$. This is the proportion of winning arguments obtained by reading...
Table 4

The results of performing a Mann–Whitney test on the dataset behind Fig. 7, in order to answer the question, “When comparing the four histograms of each policy pairwise, are they sampled from different distributions?” This is answered in the last column of the table, at a significance level of 0.05.

<table>
<thead>
<tr>
<th>Sample 1 policy</th>
<th>Sample 2 policy</th>
<th>Test statistic</th>
<th>p-value</th>
<th>Significant w.r.t. α = 0.05</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>2</td>
<td>5031068.5</td>
<td>0.0</td>
<td>Yes</td>
</tr>
<tr>
<td>1</td>
<td>3</td>
<td>3762044.0</td>
<td>0.0</td>
<td>Yes</td>
</tr>
<tr>
<td>1</td>
<td>4</td>
<td>3848325.5</td>
<td>0.0</td>
<td>Yes</td>
</tr>
<tr>
<td>1</td>
<td>3</td>
<td>7756407.5</td>
<td>(\sim 10^{-51})</td>
<td>Yes</td>
</tr>
<tr>
<td>2</td>
<td>3</td>
<td>7875127.5</td>
<td>(\sim 10^{-45})</td>
<td>Yes</td>
</tr>
<tr>
<td>3</td>
<td>4</td>
<td>9307703.0</td>
<td>0.137</td>
<td>No</td>
</tr>
</tbody>
</table>

Table 5

The summary statistics for the data behind Fig. 8.

<table>
<thead>
<tr>
<th>Policy</th>
<th>Attribute</th>
<th>Direction</th>
<th>DFS?</th>
<th>Median</th>
<th>Mean</th>
<th>SSD</th>
<th>Minimum</th>
<th>Maximum</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>Time</td>
<td>Ascending</td>
<td>No</td>
<td>0.150</td>
<td>0.172</td>
<td>0.131</td>
<td>0</td>
<td>0.571</td>
</tr>
<tr>
<td>2</td>
<td>Time</td>
<td>Ascending</td>
<td>Yes</td>
<td>0.188</td>
<td>0.203</td>
<td>0.128</td>
<td>0</td>
<td>0.600</td>
</tr>
<tr>
<td>3</td>
<td>Rational Likes</td>
<td>Descending</td>
<td>Yes</td>
<td>0.200</td>
<td>0.220</td>
<td>0.134</td>
<td>0</td>
<td>0.625</td>
</tr>
<tr>
<td>4</td>
<td>Empirical Likes</td>
<td>Descending</td>
<td>Yes</td>
<td>0.200</td>
<td>0.219</td>
<td>0.133</td>
<td>0</td>
<td>0.625</td>
</tr>
</tbody>
</table>

Table 6

The results of performing a Mann–Whitney test on the dataset behind Fig. 8, in order to answer the question, “When comparing the four histograms of each policy pairwise, are they sampled from different distributions?” This is answered in the last column of the table, at a significance level of 0.05.

<table>
<thead>
<tr>
<th>Sample 1 policy</th>
<th>Sample 2 policy</th>
<th>Test statistic</th>
<th>p-value</th>
<th>Significant w.r.t. α = 0.05</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>2</td>
<td>-12.03984</td>
<td>(\sim 10^{-33})</td>
<td>Yes</td>
</tr>
<tr>
<td>1</td>
<td>3</td>
<td>-16.97377</td>
<td>(\sim 10^{-65})</td>
<td>Yes</td>
</tr>
<tr>
<td>1</td>
<td>4</td>
<td>-16.87197</td>
<td>(\sim 10^{-64})</td>
<td>Yes</td>
</tr>
<tr>
<td>2</td>
<td>3</td>
<td>-5.460176</td>
<td>(\sim 10^{-8})</td>
<td>Yes</td>
</tr>
<tr>
<td>2</td>
<td>4</td>
<td>-5.309276</td>
<td>(\sim 10^{-8})</td>
<td>Yes</td>
</tr>
<tr>
<td>3</td>
<td>4</td>
<td>-0.175989</td>
<td>0.430</td>
<td>No</td>
</tr>
</tbody>
</table>

the first five arguments given the sorting policy, regardless of the size of the debate. This gives us the series of partially-overlapping histograms in Fig. 8.

The summary statistics of the distributions of \(J_5\) for the four policies are shown in Table 5.

When comparing medians, Policy 1 is again the least effective. Policy 2 is more effective than Policy 1 (i.e. sorting by time with DFS is better than without DFS). Both Policies 3 and 4 are more effective than Policy 2, and they seem as effective as each other. Again, we can make this precise with the same statistical tests, hypotheses and significance levels as for the summary statistics in Fig. 7 (see Table 6).

We conclude that the ordinal ranking the four policies by \(\langle J \rangle - \frac{1}{3}\) and by \(J_5\) do not change. Therefore, when seen through our two ways of measuring the effectiveness of these policies, sorting Kialo claims by likes-DFS (as defined in Section 5.2), on average, contains a higher proportion of the actually winning arguments than sorting by time (DFS or not).
6. Discussion and future work

Modern online debates are often so large in scale that the majority of Internet users cannot read all comments made. Many platforms that host these networks make use of comment sorting policies that can display the most suitable comments to readers who may not have read everything. We are interested in applying argumentation theory to measure how effective these policies are in displaying the actually winning arguments to readers who may not have read the entire network.

Our first contribution (Section 3) is a pipeline that takes as input sufficiently “clean” debates represented as bipolar argumentation trees, each having a unique set of (normatively) winning arguments. A sorting policy specifies the order to display the arguments. At each incomplete reading of the pipeline with respect to a policy, which has its own set of provisionally winning arguments from the induced sub-framework, we measure how many actually winning arguments are present using the Jaccard coefficient.

Our second contribution (Sections 4 and 5) is an application of this pipeline to Kialo debates. We argue that Kialo debates are “clean” thanks to its moderation policy. Therefore, each debate is readily-represented as a bipolar argumentation tree. The pipeline calculates, for each debate and policy, its sequence of Jaccard coefficients, one per amount of the debate read. We then aggregate these coefficients to measure how effective the policy is across all debates. As a starting point we have opted to compare policies based on sorting by likes and time, as these are the most obvious attributes available from Kialo. We find that policies that sort the comments from most to least likes are on average more effective at displaying the actually winning arguments than sorting comments from the earliest to the most recent.

This work uses Kialo as a case study on how we can apply argumentation theory to measure the effectiveness of comment sorting policies, taking advantage of Kialo’s “clean” nature. But the pipeline is agnostic to the kinds of debates and policies. In future work, we seek to determine how robust our observation that “sorting by likes is better than sorting by time” for other debates. As stated in Section 3, we cannot assume that other debates are as “clean” as Kialo, so we seek to integrate some argument mining techniques (e.g. [20], or those summarised in Section 2.3) as a pre-processing step. For example, in BBC News’ Have Your Say, the entire news article is the root argument, and the comments replying to it may not be arguments. But instead of being a star graph, some comments may mention a username or quote from another comment as a reply. Therefore, one argument mining task would be to infer replies via the text, rather than having the replies already structured like in Kialo. Another argument mining task is to infer the argumentative structure within the news article, which will then split the news article from a single root into a cluster of arguments, and one can identify the specific portions of the news text a given comment is replying to. Further, we can also investigate the time-evolution of such debates, e.g. by taking a time-ordered sequence of debate snapshots and comparing the pipeline outputs at each step.

We anticipate that for noisier datasets we may have to deal with trolls and bots that respectively make abusive comments or distribute spam (e.g. [5,18,30]). For such networks, it may be problematic to allow for the unrebutted comments to win by default, as they are the furthest away from the root and may be irrelevant or insulting. We have made an attempt at mitigating the effects of the leaf nodes winning by default [4], but other ideas involve changing the debate graph topology to have some leaves be symmetrically attacked. However, the resulting loss of a tree structure may mean our actually winning set of arguments can be non-unique, so which actually winning set of arguments (e.g. preferred or grounded) should we compare against? Or would different sorting policies guide readers to different sets of winners?

We can also consider other sorting policies. For example, in Daily Mail comments, the likes attribute contains both an upvote and a downvote. We may (e.g.) compare sorting by descending upvote with
sorting by ascending downvote. As mentioned in Section 3, Daily Mail depth-2 comments are hidden under a “See all Replies” button. How does selectively revealing such comments affect the visibility of the actually winning arguments and hence the effectiveness of a policy? The four policies we have studied are chosen because they are present in most debates (Section 3), but Kialo’s user interface (UI) also displays breadth-first search, so we can investigate how effective Kialo’s UI is as a “policy”.

As for post-processing, we have alluded to the many ways the sequence of Jaccard coefficients can be aggregated in Section 5.4. We can test the robustness of our observations by considering more fine-grained averages such as averaging over e.g. the first third of \( \frac{n}{3} \) and see how the policies fare. We may also replace the Jaccard coefficient with something more sophisticated (e.g. [26]). We may also consider more current BAF theory, specifically the various interpretations of “support”, when calculating winning arguments (e.g. [10]). Such modifications to the pipeline should distinguish whether the claim of “likes” being better than “time” is due to how the pipeline processes the data, or due to the data itself.

Finally, we seek to answer the question of why should sorting by likes be better than sorting by time. We hypothesise that likes is somehow related to how “good” a point that comment is making, which by standards of civil discourse raises its probability of winning and force of argument. Future work will make these notions precise.

Acknowledgements

The authors would like to thank the three anonymous reviewers whose suggestions have greatly improved the paper. APY and NS acknowledge funding from the Space for Sharing (S4S) project (Grant No. ES/M00354X/1). SJ was funded by a King’s India scholarship offered by the King’s College London Centre for Doctoral Studies. GB and NS acknowledge funding from the Engineering and Physical Sciences Research Council (EPSRC) through the Centre for Doctoral Training in Cross Disciplinary Approaches to Non-Equilibrium Systems (CANES, Grant No. EP/L015854/1).

References


[25] A.P. Young et al. / Ranking comment sorting policies in online debates 285